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The State of eBPF Fuzzing



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Who Am I?

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Interested in BPF fuzzing for a while:

- First contributions to verifier after rebasing bpf-fuzzer
- Updating BPF descriptions in Syzkaller since 2019

eBPF Fuzzing

- Syzkaller coverage
- Recent improvements
- Other approaches
- Finding a test oracle
- Conclusion

Syzkaller

- Well maintained and well integrated
 - Lots of up-to-date syscall descriptions
 - Continuously running and reporting bugs (syzbot)
- Code coverage guided fuzzer (kcov)
- Structured fuzzer: syzlang descriptions of syscalls
- Can find many bugs via kernel sanitizers (KASAN & co.)
- Runs full kernels in VMs

Syzkaller's Coverage

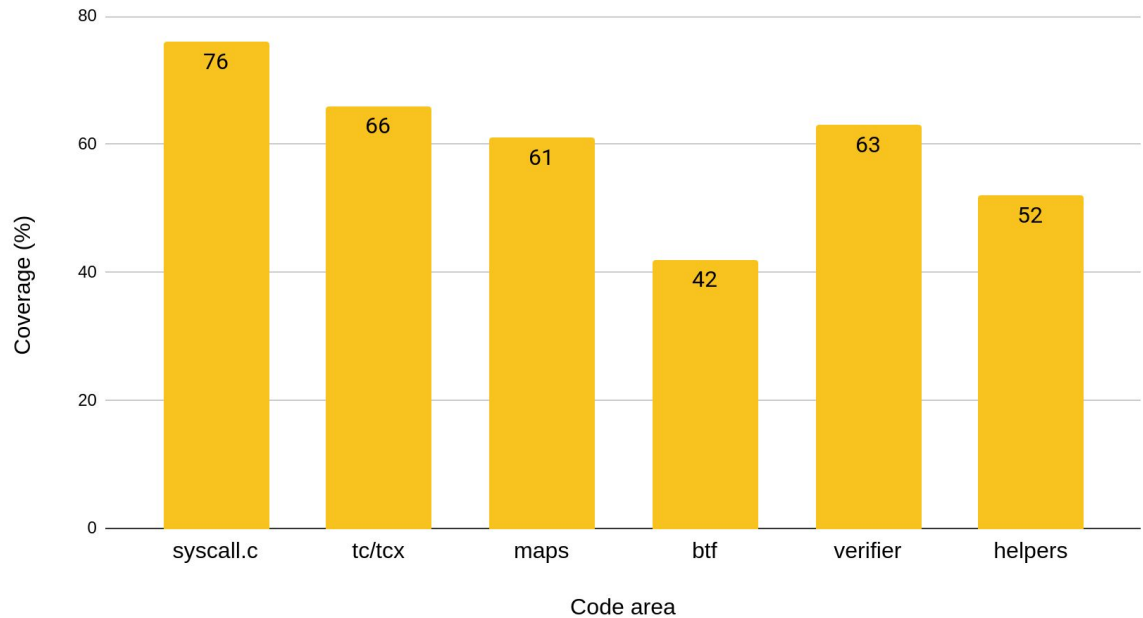
- Syzkaller finds many bugs in BPF
- But doesn't tell much on effectiveness
- Code coverage is a better measure

```
WARNING in vmap_pages_range_noflush (2) [bpf] [net]
BUG: unable to handle kernel paging request in bpf_prog_A...
WARNING in sock_map_close (2) [bpf] [net]
possible deadlock in lock_timer_base [net] [bpf]
KASAN: slab-use-after-free Read in htab_map_alloc (2) [bpf]
KASAN: stack-out-of-bounds Read in xdp_do_check_flushed...
general protection fault in __dev_flush [net] [bpf]
general protection fault in bq_flush_to_queue [bpf] [net]
KASAN: slab-use-after-free Read in bq_xmit_all [bpf] [net]
general protection fault in __xsk_map_flush [bpf] [net]
general protection fault in __cpu_map_flush [bpf] [net]
general protection fault in xdp_do_generic_redirect [net] [bpf]
general protection fault in dev_map_redirect [bpf] [net]
stack segment fault in dev_hash_map_redirect [bpf] [net]
general protection fault in xdp_do_redirect [bpf] [net]
stack segment fault in bpf_xdp_redirect [bpf] [net]
stack segment fault in cpu_map_redirect [net] [bpf]
WARNING in bpf_lwt_seg6_adjust_srh [bpf] [net]
possible deadlock in console_flush_all (2) [trace] [bpf]
WARNING in skb_ensure_writable [bpf] [net]
INFO: task hung in bpf_prog_dev_bound_destroy [bpf]
possible deadlock in __sock_map_delete [bpf] [net]
general protection fault in bpf_get_attach_cookie_tracing ...
KASAN: slab-use-after-free Read in bpf_link_free (2) [bpf]
WARNING in sock_map_close [bpf] [net]
possible deadlock in sock_hash_delete_elem (2) [bpf] [net]
```

Syzkaller's Coverage

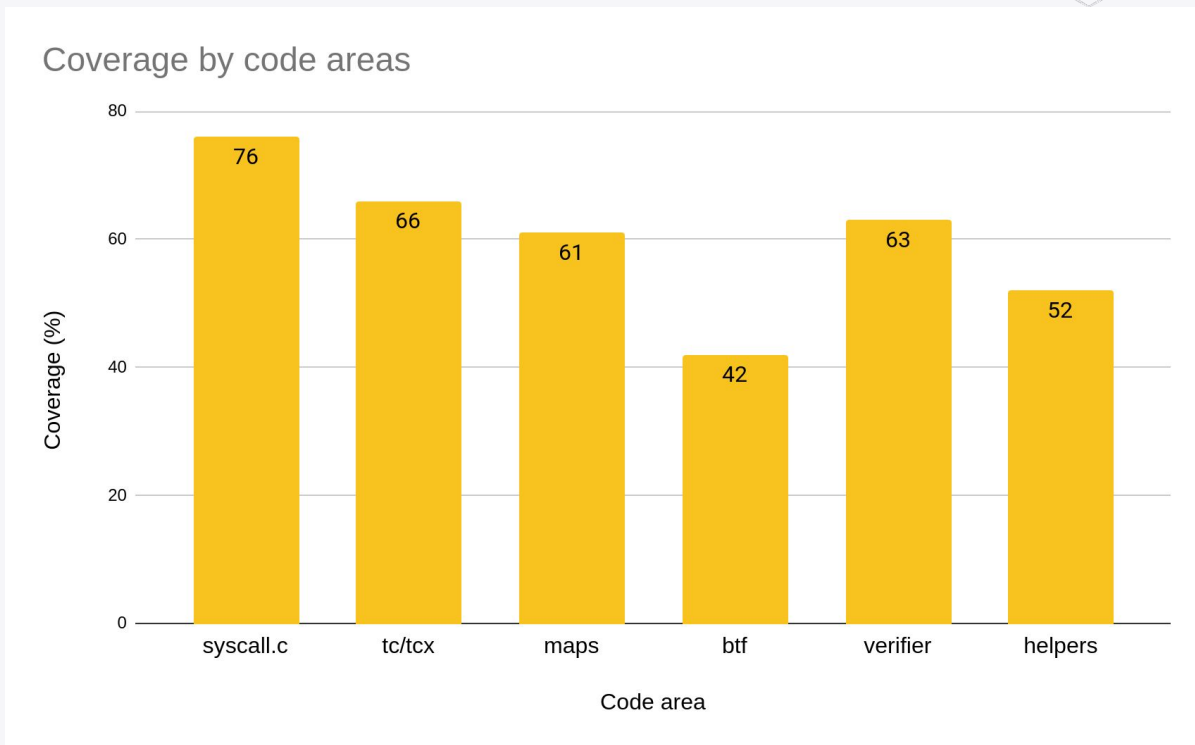
- From syzbot
- Rough aggregation

Coverage by code areas



Syzkaller's Coverage

- From syzbot
- Rough aggregation
- Closer to syscall is better
- BTF desc. is outdated
- Helpers are hard to reach



Challenges of Fuzzing BPF

- Many dependencies between various part of the input
 - Ex. program type and allowed helpers
 - Ex. sizes between map creation and map value load
 - Ex. jump offset and program structure
 - Ex. ordering between write and read of R0
 - Ex. BTF kfunc prototypes and kfunc calls
- Hard to describe with a simple description language
 - Don't want to reimplement the verifier in the fuzzer 🤪
- Several layers to pass:
 - Ex. kfunc call requires valid BTF + valid program + valid exec syscall

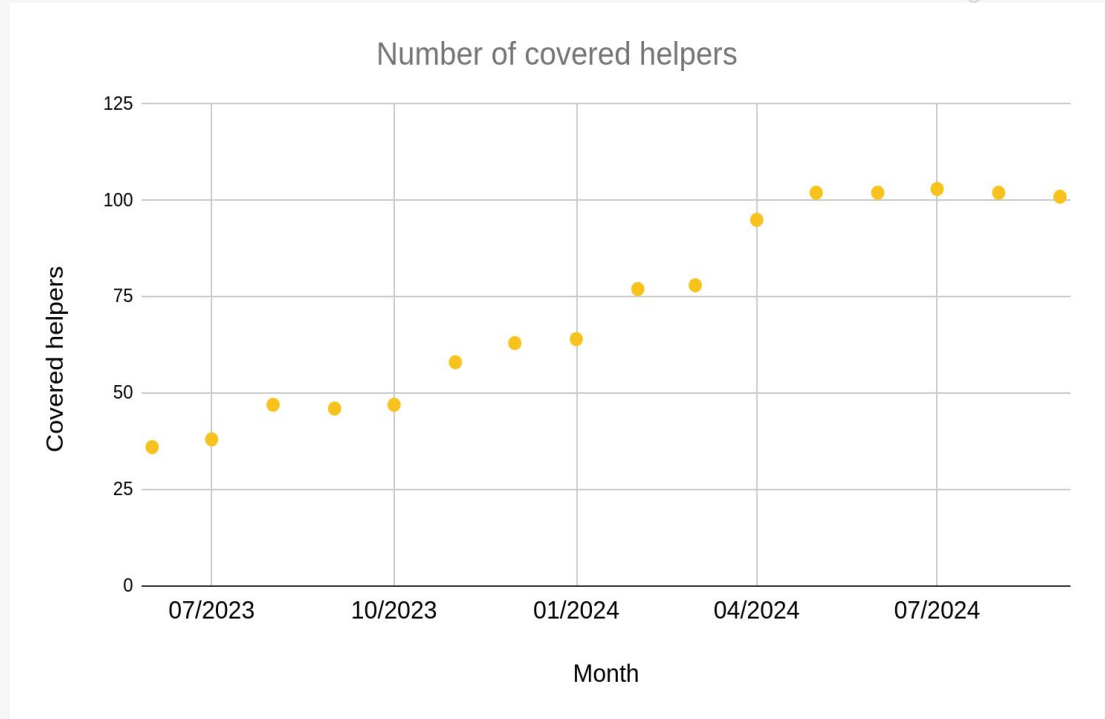


eBPF Fuzzing

- ◆ Syzkaller coverage
- ◆ **Recent improvements**
- ◆ Other approaches
- ◆ Finding a test oracle
- ◆ Conclusion

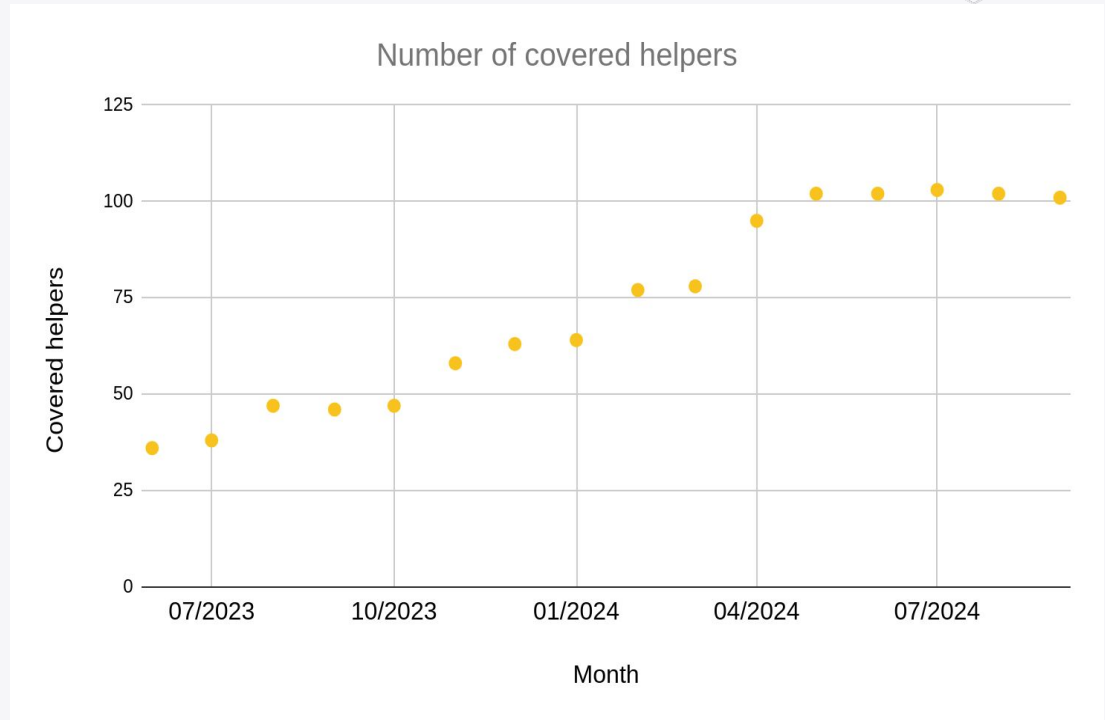
Recent Improvements

- Counting all covered helpers, even partially covered
- Covered helpers doubled in a year (+55)



Recent Improvements

- Counting all covered helpers, even partially covered
- Covered helpers doubled in a year (+55)
- Describing full helper calls paid off
- Described only 8 helpers, syzkaller guessed the rest



Recent Improvements

- Syzlang becoming more expressive with conditional fields
- Enabled more precision in BPF descriptions

```
link_create_netkit {  
    relative_link_fd      fd_bpf_link      (if[value[bpf_link_create_arg_t:flags] & BPF_F_LINK_OR_ID == BPF_F_LINK])  
    relative_prog_fd      fd_bpf_prog       (if[value[bpf_link_create_arg_t:flags] & BPF_F_LINK_OR_ID == 0])  
    relative_link_id      bpf_link_id      (if[value[bpf_link_create_arg_t:flags] & BPF_F_LINK_OR_ID == BPF_F_LINK_OR_ID])  
    relative_prog_id      bpf_prog_id      (if[value[bpf_link_create_arg_t:flags] & BPF_F_LINK_OR_ID == BPF_F_ID])  
    exp_revision          bpf_revision  
}  
 [packed]
```

eBPF Fuzzing

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- ◆ **Other approaches**
- ◆ Finding a test oracle
- ◆ Conclusion

Running the Verifier in Userspace

- bpf-fuzzer by Facebook: first-ever eBPF fuzzer
 - Uses libfuzzer
 - Verifier executed in userspace with lots of glue code
- kBdysch by Anatoly Trosinenko
 - Uses AFL
 - Relies on LKL instead of manual port to userspace
- High maintenance cost!
 - Did find multiple bugs though

Buzzer: Tailored Fuzzing for BPF

- Runs the kernel in VMs like syzkaller
- Somewhat focus on the verifier
- BPF-specific fuzzing strategies
 - Attempting out-of-bound map writes
 - Checking verifier logs
 - Or plain old coverage-based
- Found two vulnerabilities so far
- Focus of the next talk!



eBPF Fuzzing

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- ◆ **Finding a test oracle**
- ◆ Conclusion

We're Missing a Test Oracle!

- Good at finding memory errors, crashes, deadlocks, kernel warnings, etc.
- Struggle to find verifier bypasses
 - Because verifier bypasses are typically silent
- Need a test oracle for the verifier's soundness

State Embeddings as a Test Oracle

- Hao Sun and Zhendong Su devised one test oracle for the verifier
- Turn silent soundness issues into loud verifier errors
- Published at [OSDI'24](#)

State Embeddings as a Test Oracle

1. Start from accepted BPF program

```
0: *(u64*)(r10 -40) = -1
1: r1 = *(u64*)(r10 -40)
2: r2 = 1
3: if r1 < 0 goto +1
4: r2 = 0
5: exit
```

State Embeddings as a Test Oracle

1. Start from accepted BPF program
2. Fold variables into single register

```
0: r9 = 0
1: *(u64*)(r10 -40) = -1
2: r1 = *(u64*)(r10 -40)
3: r2 = 1
4: if r1 < 0 goto+1
5: r2 = 0
6: r9 += r1
7: r9 *= r2
8: if r9 != -1 goto+1
9: verifier_sink()
10: exit
```

State Embeddings as a Test Oracle

1. Start from accepted BPF program
2. Fold variables into single register
3. Trigger verifier error if folded value is as expected
 - a. Ex. write to R10

```
0: r9 = 0
1: *(u64*)(r10 -40) = -1
2: r1 = *(u64*)(r10 -40)
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State Embeddings as a Test Oracle

1. Start from accepted BPF program
2. Fold variables into single register
3. Trigger verifier error if folded value is as expected
 - a. Ex. write to R10
4. Modified program passes verifier implies:
 - a. concrete folded value \notin abstract folded value
 - b. Ex. $-1 \notin$ verifier's view of R9
 - c. That is, unsoundness issue!

```
0: r9 = 0
1: *(u64*)(r10 -40) = -1
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Conclusion

- Help welcome for syzkaller descriptions!
 - Lots to do, easy and harder stuff
 - Very helpful & responsive maintainers
 - Bugs as rewards
- IMO, we should converge approaches in syzkaller
 - Better integration with the kernel (syzbot)
- How can we improve the status quo?
 - Can we use state embeddings in syzkaller? In buzzer?

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Thanks !

